Number of quasi-kernels in digraphs

Tay Eng Guan National Institute of Education Nanyang Technological University Gregory Gutin, Khee Meng Koh, Eng Guan Tay and Anders Yeo (2004) On the number of quasi-kernels in digraphs. *Journal of Graph Theory* 46(1), 48–56. DOI: 10.1002/jgt.10169

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A vertex set X of a digraph D = (V, A) is a kernel if X is independent (i.e., all pairs of distinct vertices of X are non-adjacent) and for every $v \in V - X$ there exists $x \in X$ such that $\underline{vx} \in A$. A vertices of A are non-adjacent) and for every $v \in V$ -X there exists $x \in X$ such that $\underline{w} \in A$. A vertex set X of a digraph D = (V, A), is a quasi-kernel if X is independent and for every $v \in V$, there exist $w \in V$ -X, $x \in X$ such that either $\underline{w} \in A$ or $\underline{v}\underline{w}$, $\underline{w} \in A$. In 1974, Chyátal and Lovász proved that every digraph has a quasi-kernel. In 1996, Jacob and Meyniel proved that if a digraph D has no kernel, then D contains at least three quasi-kernels. We characterize digrapl with exactly one and two quasi-kernels, and, thus, provide necessary and sufficient condition adjoined by boxes of lose three gains in kernels. a digraph to have at least three quasi-kernels. In particular, we prove that every strong digrap order at least three, which is not a 4-cycle, has at least three quasi-kernels.

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Definition 1

A vertex set X of a digraph D = (V, A) is a kernel if X is independent and for every $v \in V - X$, there exists $x \in X$, such that $vx \in A$.

Definition 2

A vertex set X of a digraph D = (V, A) is a quasi-kernel if X is independent and for every $v \in V - X$, there exist $w \in V - X$, $x \in X$, such that either $vx \in A$ or vw, $vx \in A$

Theorem (Chvátal & Lovász, 1974)

Every digraph has a quasi-kernel.

Theorem (Jacob & Meyniel, 1996)

If a digraph D has no kernel, then D has at least 3 quasikernels.

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Lemma 2.1 Let x be a vertex in a digraph D. If x is a non-sink, then D has a quasi-kernel not including x.

Proof: Let $y \in N^+[x] - \{x\}$ be arbitrary. If $N^-[y] = V(D)$, then y is the required quasi-kernel. If $N^-[y] \neq V(D)$, let Q' be a quasi-kernel in $D - N^-[y]$. If y does not a vertex in Q', then Q' is a quasi-kernel in D, which does not contain x. If y does not dominate a vertex in Q', then $Q' \cup \{y\}$ is a quasi-kernel in D, which does not include $x.\square$

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Theorem 2.2 A digraph D has only one quasi-kernel if and only if D has a sink and every non-sink of D dominates a sink of D. If a digraph D has only one quasi-kernel Q, then Q is a kernel and consists of the sinks of D.

Proof: Assume that D has a sink and every non-sink of D dominates a sink of D. Let S be the set of sinks in D. To see that S is a unique quasi-kernel of D, it is enough to observe that every sink must be in a quasi-kernel.

Let D have only one quasi-kernel Q. To see that Q is the set of sinks in D, observe that Q contains all sinks in D and, by Lemma 2.1, Q does not have non-sinks. If x is a non-sink and x does not dominate a vertex in Q, then $Q \cup \{x\}$ is another quasi-kernel of D, a contradiction. Thus, we have proved that D has a sink and every non-sink of D.

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In view of Theorem 2.2, the following assertion is a strengthening of the Jacob-Meyniel theorem for the case of digraphs with no sinks.

Theorem 2.3 Let D be a digraph with no sink. Then D has precisely two quasi-kernels if and only if D has an induced 4-cycle or 2-cycle, C, such that no vertex of C dominates a vertex in D-V(C) and every vertex in D-V(C) dominates at least two adjacent vertices in C.

Theorem 2.5 A strong digraph D of order at least three has at least three quasi-kernels, unless D is \vec{C}_4 .

Proof: Immediate from the previous theorems, Theorems 2.2 and 2.3.

Theorem 2.6 Let D be a digraph, S the set of sinks in D, R the set of vertices that have an arc into S, and H = D - S - R. Then D has precisely two quasi-kernels, if and only if one of the following holds:

- (a) There is a 2-cycle C in H such that at most one of the vertices in C has an arc into R, no vertex of C dominates a vertex in H-V(C), and every vertex in H-V(C) dominates both vertices in C.
- (b) There is an induced 4-cycle, C, in H such that no vertex of C dominates a vertex in D-V(C) and every vertex in H-V(C) dominates two adjacent vertices in C.
- (c) The digraph H has at least two vertices. There is a vertex x in H such that no vertex of H is dominated by x, all the vertices of H-x dominate x, i.e., $(V(H)-\{x\},x)=(V(H)-\{x\})\times\{x\}$, and there is a kernel Q in H-x, consisting only of sinks in H-x. Moreover, there is no arc from Q to R.
 - (d) The digraph H has exactly one vertex and this vertex dominates a vertex in R.

3 Disjoint quasi-kernels

If a digraph D has a sink x, then every quasi-kernel in D must contain x. Hence, a digraph with sinks has no disjoint quasi-kernels. However, one may suspect that every digraph with no sink has a pair of disjoint quasi-kernels. By Lemma 2.1, this is true for digraphs with exactly two quasi-kernels: see the first paragraph in the proof of Theorem 2.3. One can show that this is also true for every digraph which possesses a quasi-kernel of cardinality at most two.

Unfortunately, in general, the above claim does not hold. Consider the following construction suggested to us by the referee. Let T be a tournament having the property that for every pair x, y of vertices there exists a vertex z such that $x \rightarrow z$ and $y \rightarrow z$. The existence of such tournaments was first proved by Erdős [4], see also Section 1.2 in [1]. It was shown by Graham and Spencer [5] that some quadratic residue tournaments are such tournaments, see also Section 9.1 in [1].) Extend T to a digraph D by adding, for every vertex x in T, a new vertex x together with the arc x'x.

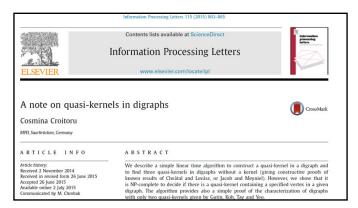
Clearly, D has no sink and every quasi-kernel of D contains exactly one vertex in T. If Q_x and Q_y are a pair of quasi-kernels of D containing the vertices x and y, respectively, then they are not disjoint because they both have to contain z', where $x \rightarrow z$ and $y \rightarrow z$.

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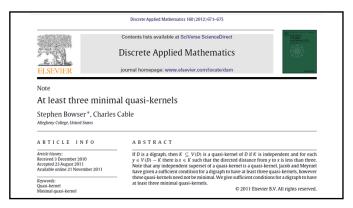
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Scott Heard & Jing Huang (2008). Disjoint quasi-kernels in digraphs

Abstract: A quasi-kernel in a digraph is an independent set of vertices such that any vertex in the digraph can reach some vertex in the set via a directed path of length at most two. Chyátal and Lovász proved that every digraph has a quasi-kernel. Recently, Gutin et al. raised the question of which digraphs have a pair of disjoint quasi-kernels. Clearly, a digraph has a pair of disjoint quasi-kernels cannot contain sinks, that is, vertices of out-degree zero, as each such vertex is necessarily included in a quasi-kernel. However, there exist digraphs which contain neither sinks nor a pair of disjoint quasi-kernels. Thus, containing no sinks is not sufficient in general for a digraph to have a pair of disjoint quasi-kernels. In contrast, we prove that, for several classes of digraphs, the condition of containing no sinks guarantees the existence of a pair of disjoint quasi-kernels. The classes contain semicomplete multipartite, quasi-transitive, and locally semicomplete digraphs. © 2008 Wiley Periodiculs, Inc. J Graph Theory 58: 251–260, 2008



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A Result on the Small Quasi-Kernel Conjecture Allan van Hulst VU University Amsterdam
The Netherlands
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Any directed graph D=V(D), A(D) in this work is assumed to be finite and without sile-loops. A source in a directed graph is a vertex having at least one integring are, A quasirent of $\mathbb{C}\times V(D)$ is an independent set in D when that every vertex in V(D) can be readed in a most two steps from a vertex in V(D) can be readed in a most two steps from a vertex in V(D) is as open V(D) and V(D) can be readed in a most two steps from a vertex in V(D) and of this paper is to prove the SQRC under the assumption of a structural property of directed graphs. The sins of this paper is to prove the SQRC under the assumption of a structural property of directed graphs. The sinst the SQRC to the existence of a vertex V(D) and a bound on the number of new sources emerging when v and it is out-neighborhood and direction V(D) and V(D) in the single contraction V(D) and V(D) in the single contraction V(D) and V(D) in the single contraction V(D) in the single contraction V(D) in V(D) is a first V(D) and a bound on the number of new sources emerging when V(D) in V(D) in V(D) in V(D) is a first V(D) and V(D) in V(

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